Profitable entanglement for channel discrimination

Samad Khabbazi Oskouei,^{1, *} Stefano Mancini,^{2,3,¶} and Milajiguli Rexiti^{2,3, ‡}

¹Department of Mathematics, Varamin-Pishva Branch Islamic Azad University, Varamin, 33817-7489 Iran

²School of Science and Technology, University of Camerino, Via Madonna delle Carceri 9, I-62032 Camerino, Italy NEN, Spinso Perusia, Via A, Passeli, 10(122 Perusia, Ital)

³INFN–Sezione Perugia, Via A. Pascoli, I-06123 Perugia, Italy

We investigate the usefulness of side entanglement in discriminating between two generic qubit channels, up to unitary pre- and post-processing, and determine exact conditions under which it does enhance (as well as conditions under which it does not) the success probability. This is done in a constructive way by first analyzing the problem for channels that are extremal in the set of completely positive and trace-preserving qubit linear maps and then for channels that are inside such a set.

Keywords: quantum channels, entanglement, statistical distance.

I. INTRODUCTION

The relevance of discerning between two (or more) quantum channels stems from the fact that any physical process can be described as a quantum channel [1–4]. It becomes even more pronounced when quantum channels are employed to encode information, like in quantum reading [5]. Broadly speaking, the task of channel discrimination is challenging [6, 7]. In fact, albeit for a given input it can be traced back to the problem of state discrimination [8], it involves optimization of overall inputs. The performance is usually expressed in terms of minimum error probability. Entanglement offers the possibility of enhancing such a performance by resorting to a strategy that uses a pair of systems such that only one undergoes the effect of the channel while both are measured (a strategy usually referred to as *side entanglement*)[20]. Remarkably this strategy can also enhance the distinguishability of entanglement-breaking channels [10]. It is known that an ancilla with the same dimension as the input of the channels is always sufficient for optimal discrimination, while this is not the case for the output dimension whenever it is smaller [11].

However, until now, there is no clear and comprehensive picture of which channels can benefit from side entanglement, even at the level of the qubit. Here we clarify this point by resorting to the characterization of qubit channels in terms of affine maps on \mathbb{R}^3 [12].

We find conditions under which side entanglement does enhance (as well as conditions under which it does not) the performance in discriminating between two generic qubit channels, up to unitary pre- and post-processing. This is achieved in a constructive way by first analyzing the problem for channels that are extremal in the set \Re of completely positive and trace-preserving qubit linear maps, up to unitary pre- and post-processing, and then for channels that are inside such a set.

II. PRELIMINARIES

A quantum channel is a linear, completely positive and trace preserving (CPTP) map on the set $\mathcal{D}(\mathcal{H})$ of density operators over a Hilbert space \mathcal{H} . Every quantum channel $\mathcal{M} : \mathcal{D}(\mathcal{H}) \to \mathcal{D}(\mathcal{H})$

^{*} kh_oskuei@yahoo.com

[¶] stefano.mancini@unicam.it

⁴ milajiguli.milajiguli@gmail.com

can be expressed in the operator sum (Kraus) representation:

$$\mathcal{M}(\rho) = \sum_{j} K_{j} \rho K_{j}^{\dagger}, \tag{1}$$

where $K_j : \mathcal{H} \to \mathcal{H}$ are linear operators (called Kraus operators) satisfying the normalization condition $\sum_j K_j^{\dagger} K_j = I$. The number of non-zero operators in the Kraus representation is called the Kraus rank.

In the following we will consider only qubit channels. For them, it is known that the maximum Kraus rank is four.

In such a case, given a density operator $\rho = \frac{1}{2}(I + \mathbf{r} \cdot \boldsymbol{\sigma})$ with $\mathbf{r} \in \mathbb{R}^3$, $\|\mathbf{r}\| \leq 1$, and $\boldsymbol{\sigma}$ the vector of Pauli operators, the channel \mathcal{M} can always be represented as an affine transformation in \mathbb{R}^3 , that is

$$r \xrightarrow{\mathcal{M}} M r + t,$$
 (2)

where M is a 3×3 real matrix and $t = (t_1, t_2, t_3)^T \in \mathbb{R}^3$.

The matrix M, via orthogonal transformations, can be brought to diagonal form $N = diag(\lambda_1, \lambda_2, \lambda_3)$. By virtue of that, any qubit channel M can be written as

$$\mathcal{M}(\bullet) = V\mathcal{N}\left(U \bullet U^{\dagger}\right)V^{\dagger},\tag{3}$$

where N is a channel with associated diagonal matrix N and vector t, while U, V are suitable (pre- and post-processing) quibit unitaries.

Then, the image through N of the boundary of the Bloch sphere (pure states) is the ellipsoid described by the equation

$$\left(\frac{x-t_1}{\lambda_1}\right)^2 + \left(\frac{y-t_2}{\lambda_2}\right)^2 + \left(\frac{z-t_3}{\lambda_3}\right)^2 = 1,$$
(4)

where λ_k s define the length of the axes of the ellipsoid and *t* its center.

Note that not all ellipsoids within the Bloch sphere correspond to a quantum channel (completely positive map).

Now let us consider the set of qubit channels

$$\mathfrak{N} := \left\{ \mathcal{N} : \mathcal{D}(\mathbb{C}^2) \to \mathcal{D}(\mathbb{C}^2) \,|\, \mathcal{N} \text{ is linear and CPTP and } \mathcal{N} \leftrightarrow (\mathbf{N}, t) \right\}.$$
(5)

When N belongs to the closure of extreme points of \mathfrak{N} , it can be parametrized by (see Sec. II.1 of Ref. [13] and reference therein)

$$\boldsymbol{t} = \begin{pmatrix} 0 \\ 0 \\ \sin(\phi - \theta)\sin(\phi + \theta) \end{pmatrix}, \quad \boldsymbol{N} = \begin{pmatrix} \cos(\phi - \theta) & 0 & 0 \\ 0 & \cos(\phi + \theta) & 0 \\ 0 & 0 & \cos(\phi - \theta)\cos(\phi + \theta) \end{pmatrix}, \quad (6)$$

with $\theta, \phi \in [0, \pi]$. In term of Kraus operators we have

$$K_0 = \begin{pmatrix} \cos \theta & 0 \\ 0 & \cos \phi \end{pmatrix}, \quad K_1 = \begin{pmatrix} 0 & \sin \phi \\ \sin \theta & 0 \end{pmatrix}.$$
(7)

When t = 0 (unital channels) the extreme points of \mathfrak{N} are known to be the maps that conjugate by a unitary matrix and in the λ_k representation they correspond to four corners of a tetrahedron. The channels on the edges of such a tetrahedron correspond to ellipsoids that have exactly two points in common with the boundary of the Bloch sphere. These are called *quasi-extreme points* because looking at extremality within \mathfrak{N} they belong to the closure of extreme points of \mathfrak{N} , but are not true extreme points.

When $t \neq 0$ (non-unital channels) the extreme points of \mathfrak{N} correspond to maps for which a translation allows the ellipsoid to touch the boundary of the Bloch sphere at two points.[21]

It is known [12] that any quantum channel on \mathfrak{N} can be written as convex combination of two maps in the closure of extreme points. In particular, a unital channel can be written as a convex combination of two maps on the edges of the tetrahedron.

A. Channel Discrimination

Let \mathcal{N}_{ω_1} and \mathcal{N}_{ω_2} be two qubit channels characterized respectively by parameters $\omega_1, \omega_2 \in \Omega$.[22]

Suppose each of the above channels appears with probability $\frac{1}{2}$, and we want to discriminate between them by inputting a probe state $\rho \in \mathcal{D}(\mathbb{C}^2)$ and measuring the output. The problem can be traced back to quantum states discrimination. According to [8], given two quantum states $\mathcal{N}_{\omega_0}(\rho)$ and $\mathcal{N}_{\omega_1}(\rho)$ appearing with equal probability, the maximum probability of success in discriminating between them reads

$$P_{succ}(\rho) = \frac{1}{2} \left(1 + \frac{1}{2} \| \Delta_{\mathcal{N}}(\rho) \|_{1} \right),$$
(8)

where

$$\Delta_{\mathcal{N}}(\rho) := \mathcal{N}_{\omega_1}(\rho) - \mathcal{N}_{\omega_2}(\rho), \qquad (9)$$

and $\|\cdot\|_1$ denotes the trace-norm of an operator.

To achieve (8), the optimal measurement turns out to be the projection onto positive and negative eigenspaces of the operator $\Delta_N(\rho)$.

Then, the problem of channel discrimination becomes finding the optimal input state ρ which maximizes the trace distance of the output states.

This optimization can be restricted to the set of pure states $\mathcal{P}(\mathbb{C}^2)$ thanks to the convexity of the trace norm.

B. Side Entanglement

When a reference system is accessible for measurement, entanglement across reference and main systems can be exploited for the purpose of channel discrimination. We refer to this strategy as channel discrimination with *side entanglement*.

In this case we have to consider

$$\Delta_{\mathrm{id}\otimes\mathcal{N}}(\rho) := (\mathrm{id}\otimes\mathcal{N}_{\omega_1})(\rho) - (\mathrm{id}\otimes\mathcal{N}_{\omega_2})(\rho), \tag{10}$$

where $\rho \in \mathcal{P}(\mathbb{C}^2 \otimes \mathbb{C}^2)$, assuming without loss of generality the reference system isomorphic to the main one.

According to Eq.(8), side entanglement turns out to be profitable if there exists a parameter's region (subset of Ω) where the following inequality holds true:

$$\sup_{\rho \in \mathcal{P}(\mathbb{C}^2 \otimes \mathbb{C}^2)} \|\Delta_{\mathrm{id} \otimes \mathcal{N}}(\rho)\|_1 > \sup_{\rho \in \mathcal{P}(\mathbb{C}^2)} \|\Delta_{\mathcal{N}}(\rho)\|_1.$$
(11)

Since we will always deal with optimization on the set of pure states, be it $\mathcal{P}(\mathbb{C}^2 \otimes \mathbb{C}^2)$ or $\mathcal{P}(\mathbb{C}^2)$, in the following we will use interchangeably $|\rho\rangle\langle\rho|$ and ρ .

In the following we shall consider the discrimination between two channels \mathcal{N}_{ω_1} and \mathcal{N}_{ω_2} belonging to the set \mathfrak{N} . This is equivalent to the discrimination between generic channels with associated matrices M_1 and M_2 that can be simultaneously diagonalized (or in other words two channels that have the same pre- and post-processing unitaries).

III. DISCRIMINATING BETWEEN EXTREMAL QUBIT CHANNELS

In this Section we start considering the discrimination between two quasi-extremal qubit channels and then move on to the discrimination between two extremal qubit channels.

Theorem III.1 *In the discrimination between two qubit channels that are quasi-extreme point maps of* \mathfrak{N} *, side entanglement is not useful.*

Proof According to the parametrization (6) for quasi-extreme point maps we must have $\sin \theta = \sin \phi$. This can happen with $\cos \theta = \cos \phi$ or with $\cos \theta = -\cos \phi$. In the first case, by (7), states along the *x* direction of \mathbb{R}^3 are left invariant. It follows that the optimal input state for channel discrimination (namely, the one that is mostly affected by \mathcal{N}_{ω}) lies in the y - z plane. Hence, it can be written as $|\psi\rangle = \sqrt{r_0} |+y\rangle + \sqrt{r_1} |-y\rangle$, where $|\pm y\rangle$ denote the eigenvectors of Pauli operator *Y*. As a consequence we have to optimize over r_0, r_1 , such that $r_0 + r_1 = 1$.

Now consider a two-qubit entangled state $|\Psi\rangle$. By the Schmidt decomposition $|\Psi\rangle = \sum_{j=\pm} \sqrt{t_j} |\tau_j\rangle |\tau_j\rangle$, where $\{|\tau_j\rangle\}_j$ are orthonormal and $t_+ + t_- = 1$. However, by local unitaries that do not affect the entanglement of $|\Psi\rangle$ we can rewrite it as $|\Psi\rangle = \sum_{j=\pm} \sqrt{\lambda_j} |jy\rangle |jy\rangle$. This means that we can express $\Delta_{id\otimes\mathcal{N}}(\Psi)$ as a 2 × 2 matrix restricting on the subspace spanned by $\{|+y\rangle |+y\rangle, |-y\rangle |-y\rangle\}$ and optimize over λ_+, λ_- satisfying $\lambda_+ + \lambda_- = 1$. However, this exactly coincides with the optimization over r_0, r_1 of the state $|\psi\rangle$ without entanglement.

In case $\cos \theta = -\cos \phi$, instead, by (7) we can realize that the states $|\pm y\rangle$ are flipped one into another. Hence they will be the most affected by the channel's action. Then, repeating the above reasoning for the entangled state $|\Psi\rangle$ we end up with the fact that $\|\Delta_{id\otimes\mathcal{N}}(\Psi)\|_1 \leq \|\Delta_{\mathcal{N}}(|\pm y\rangle\langle\pm y|)\|_1$.

Remark The channels of Theorem III.1 are on the edges of the tetrahedron and correspond to ellipsoids that have exactly two points in common with the boundary of the Bloch sphere. Examples of this kind of channels are those having two Kraus operators, one proportional to the identity and one proportional to one Pauli operator (bit-flip, phase-flip – see also [14]–, bit-phase-flip). Also those having two Kraus operators proportional to two distinct Pauli operators.

Let us now address the discrimination of qubit channels that are extreme point in \mathfrak{N} . We denote any such a channel as $\mathcal{N}_{\phi,\theta}$, since it is characterized by Kraus operators (7) with parameters ϕ, θ .

Consider the following states:

$$|\psi\rangle = \sum_{i=0}^{1} a_i |i\rangle, \qquad (12)$$

$$|\Psi\rangle = \sum_{i,j=0}^{1} a_{ij} |ij\rangle, \qquad (13)$$

with $a_i, a_{ij} \in \mathbb{C}$ such that $\sum_{i=0}^{1} |a_i|^2 = 1$ and $\sum_{i,j=0}^{1} |a_{ij}|^2 = 1$ respectively.

We want to study $\|\Delta_{\mathcal{N}_{\phi,\theta}}(\Psi)(\psi)\|_1$ vs $\|\Delta_{\operatorname{id}\otimes\mathcal{N}_{\phi,\theta}}(\Psi)\|_1$ for the channel $\mathcal{N}_{\phi,\theta}$. Define

$$\alpha := \cos^2 \theta_1 - \cos^2 \theta_2, \tag{14}$$

$$\beta \coloneqq \cos^2 \phi_1 - \cos^2 \phi_2,\tag{15}$$

$$\gamma_1 := \cos \phi_1 \cos \theta_1 - \cos \phi_2 \cos \theta_2, \tag{16}$$

$$\gamma_2 := \sin \phi_1 \sin \theta_1 - \sin \phi_2 \sin \theta_2, \tag{17}$$

together with

$$\gamma_m := \begin{cases} \gamma_1, & |\gamma_1| \le |\gamma_2| \\ \gamma_2, & |\gamma_1| > |\gamma_2| \end{cases}, \qquad \gamma_M := \begin{cases} \gamma_1, & |\gamma_1| \ge |\gamma_2| \\ \gamma_2, & |\gamma_1| < |\gamma_2| \end{cases},$$
(18)

and

$$P := \begin{cases} \alpha, & |\alpha| \ge |\beta| \\ \beta, & |\alpha| < |\beta| \end{cases}.$$
(19)

Lemma III.2 For a qubit channel $\mathcal{N}_{\phi,\theta}$ that is an extreme point in \mathfrak{N} , it is

$$\max_{\psi \in \mathcal{P}(\mathbb{C}^2)} \|\Delta_{\mathcal{N}_{\phi,\theta}}(\psi)\|_1 = \max_{0 \le s \le 1} 2 \sqrt{((1-s)\alpha + s\beta)^2 + 4s(1-s)\left[\left(\frac{|\gamma_1| + |\gamma_2|}{2}\right)^2 - \alpha\beta\right]},$$
(20)

$$= \begin{cases} \frac{(|\gamma_1|+|\gamma_2|)\sqrt{(|\gamma_1|+|\gamma_2|)^2 - 4\alpha\beta}}{\sqrt{(|\gamma_1|+|\gamma_2|)^2 - (\alpha+\beta)^2}}, & \text{if } |\alpha+\beta| < |\gamma_1|+|\gamma_2|\\ \\ 2|P|, & \text{if } |\alpha+\beta| \ge |\gamma_1|+|\gamma_2| \end{cases}.$$

Proof With the input state (12), taking into account (7), we have

$$\mathcal{N}_{\phi,\theta}(\psi) = K_0 |\psi\rangle \langle \psi | K_0^{\dagger} + K_1 |\psi\rangle \langle \psi | K_1^{\dagger} = |\psi'\rangle \langle \psi' | + |\psi''\rangle \langle \psi'' |, \qquad (21)$$

where

$$|\psi'\rangle = a_0 \cos\theta |0\rangle + a_1 \cos\phi |1\rangle, \quad |\psi''\rangle = a_1 \sin\phi |0\rangle + a_0 \sin\theta |1\rangle.$$
(22)

Then, according to Eq.(9), we will have:

$$\Delta_{\mathcal{N}_{\phi,\theta}}(\psi) = \Delta_{\psi'} + \Delta_{\psi''},\tag{23}$$

with the quantities defined, in the canonical basis, as:

$$\Delta_{\psi'} := \begin{pmatrix} |a_0|^2 \alpha & a_0 a_1^* \gamma_1 \\ a_0^* a_1 \gamma_1 & |a_1|^2 \beta \end{pmatrix}, \qquad \Delta_{\psi''} = \begin{pmatrix} -|a_1|^2 \beta & a_0^* a_1 \gamma_2 \\ a_0 a_1^* \gamma_2 & -|a_0|^2 \alpha \end{pmatrix}.$$
 (24)

It follows that

$$\|\Delta_{\mathcal{N}_{\phi,\theta}}(\psi)\|_{1} = 2\sqrt{(|a_{0}|^{2}\alpha - |a_{1}|^{2}\beta)^{2} + |a_{0}a_{1}^{*}\gamma_{1} + a_{0}^{*}a_{1}\gamma_{2}|^{2}}.$$
(25)

Without loss of generality, we assume $a_0 \in \mathbb{R}$ and $a_1 = |a_1|e^{i\eta} \in \mathbb{C}$. Taking into account the normalization condition, and the fact that $|\gamma_1 + e^{2i\eta}\gamma_2|^2$ reaches the maximum $(|\gamma_1| + |\gamma_2|)^2$, we are left with finding

$$\max_{0 \le |a_1| \le 1} 2\sqrt{(\alpha - (\alpha + \beta)|a_1|^2)^2 + |a_1|^2(1 - |a_1|^2)(|\gamma_1| + |\gamma_2|)^2},$$
(26)

giving Eq.(20) for

$$\begin{cases} |a_1|^2 = \frac{2\alpha(\alpha+\beta) - (|\gamma_1|+|\gamma_2|)^2}{2(\alpha+\beta)^2 - 2(|\gamma_1|+|\gamma_2|)^2}, & \text{if } |\alpha+\beta| < |\gamma_1|+|\gamma_2| \\ |a_1|^2 = 0, & \text{if } |\alpha+\beta| \ge |\gamma_1|+|\gamma_2| \text{ and } |\alpha| \ge |\beta| \\ |a_1|^2 = 1, & \text{if } |\alpha+\beta| \ge |\gamma_1|+|\gamma_2| \text{ and } |\alpha| \le |\beta| \end{cases}$$

$$(27)$$

Lemma III.3 For a qubit channel $\mathcal{N}_{\phi,\theta}$ that is an extreme point in \mathfrak{N} , in finding $\max_{\Psi \in \mathcal{P}(\mathbb{C}^2 \otimes \mathbb{C}^2)} \|\Delta_{\mathrm{id} \otimes \mathcal{N}_{\phi,\theta}}(\Psi)\|_1$ it is enough to search in the two dimensional subspace $span\{|00\rangle, |11\rangle\}$ or $span\{|01\rangle, |10\rangle\}$, and it results

$$\max_{\Psi \in \mathcal{P}(\mathbb{C}^2 \otimes \mathbb{C}^2)} \|\Delta_{\mathrm{id} \otimes \mathcal{N}_{\phi,\theta}}(\Psi)\|_1 = \max_{0 \le s \le 1} \frac{1}{2} \sum_{i,j=1}^2 \left| (1-s)\alpha + s\beta + (-1)^i \sqrt{((1-s)\alpha + s\beta)^2 + 4s(1-s)(\gamma_j^2 - \alpha\beta)} \right|.$$
(28)

Proof First, we find an upper bound for $\|\Delta_{id \otimes N_{\phi,\theta}}(\Psi)\|_1$, where $|\Psi\rangle$ is given in Eq.(13):

$$\begin{split} \|\Delta_{\mathrm{id}\otimes\mathcal{N}_{\phi,\theta}}(\Psi)\|_{1} &= \left\| \mathrm{id}\otimes\left(K_{0}^{(1)}K_{0}^{(1)\dagger} - K_{0}^{(2)}K_{0}^{(2)\dagger}\right)\Psi + \mathrm{id}\otimes\left(K_{1}^{(1)}K_{1}^{(1)\dagger} - K_{1}^{(2)}K_{1}^{(2)\dagger}\right)\Psi \right\|_{1} \\ &\leq \left\| \mathrm{id}\otimes\left(K_{0}^{(1)}K_{0}^{(1)\dagger} - K_{0}^{(2)}K_{0}^{(2)\dagger}\right)\Psi \right\|_{1} + \left\| \mathrm{id}\otimes\left(K_{1}^{(1)}K_{1}^{(1)\dagger} - K_{1}^{(2)}K_{1}^{(2)\dagger}\right)\Psi \right\|_{1} \\ &= \|C\|_{1} + \|D\|_{1}, \end{split}$$

$$(29)$$

where the superscripts refer to the two channels and the matrices C and D are defined as

$$C := \begin{pmatrix} \alpha_1 |a_{00}|^2 & \gamma_1 a_{00} a_{01}^* & \alpha_1 a_{00} a_{01}^* & \gamma_1 a_{00} a_{11}^* \\ \gamma_1 a_{00}^* a_{01} & \beta_1 |a_{01}|^2 & \gamma_1 a_{01} a_{10}^* & \beta_1 a_{01} a_{11}^* \\ \alpha_1 a_{00}^* a_{10} & \gamma_1 a_{01}^* a_{10} & \alpha_1 |a_{10}|^2 & \gamma_1 a_{10} a_{11}^* \\ \gamma_1 a_{00}^* a_{11} & \beta_1 a_{01}^* a_{11} & \gamma_1 a_{10}^* a_{11} & \beta_1 |a_{11}|^2 \end{pmatrix},$$

$$(30)$$

$$D := \begin{pmatrix} \alpha_2 |a_{01}|^2 & \gamma_2 a_{00}^* a_{01} & \alpha_2 a_{01} a_{11}^* & \gamma_2 a_{01} a_{10}^* \\ \gamma_2 a_{00} a_{01}^* & \beta_2 |a_{00}|^2 & \gamma_2 a_{00} a_{11}^* & \alpha_2 a_{00} a_{10}^* \\ \alpha_2 a_{01}^* a_{11} & \gamma_2 a_{00}^* a_{11} & \alpha_2 |a_{11}|^2 & \gamma_2 a_{10}^* a_{11} \\ \gamma_2 a_{01}^* a_{10} & \alpha_2 a_{00}^* a_{10} & \gamma_2 a_{10} a_{11}^* & \beta_2 |a_{10}|^2 \end{pmatrix}.$$

$$(31)$$

The non-zero eigenvalues of C and D are respectively:

$$\frac{1}{2} \left(\alpha (|a_{00}|^2 + |a_{10}|^2) + \beta (|a_{11}|^2 + |a_{01}|^2) \right) \\
\pm \frac{1}{2} \sqrt{\left((|a_{00}|^2 + |a_{10}|^2)\alpha - (|a_{01}|^2 + |a_{11}|^2)\beta \right)^2 + 4\gamma_1^2 (|a_{00}|^2 + |a_{10}|^2) (|a_{01}|^2 + |a_{11}|^2)}, \quad (32)$$

and

$$\frac{1}{2} \left(-\beta (|a_{11}|^2 + |a_{01}|^2) - \alpha (|a_{00}|^2 + |a_{10}|^2) \right) \\ \pm \frac{1}{2} \sqrt{\left(-(|a_{00}|^2 + |a_{10}|^2)\beta + (|a_{01}|^2 + |a_{11}|^2)\alpha \right)^2 + 4\gamma_2^2 (|a_{00}|^2 + |a_{10}|^2) (|a_{01}|^2 + |a_{11}|^2)}.$$
(33)

Next, we show that the subspace we have to consider for the optimization must be either $span\{|00\rangle, |11\rangle\}$, or $span\{|01\rangle, |10\rangle\}$. This is done by showing that in such cases the upper bound (29) is actually achievable.

Let us take $|\Phi\rangle = a_0 |00\rangle + a_1 |11\rangle$, with $|a_0|^2 + |a_1|^2 = 1$, then we can find that

$$\mathrm{id} \otimes \mathcal{N}_{\phi,\theta}(\Phi) = K_0 |\Phi\rangle \langle \Phi| K_0^{\dagger} + K_1 |\Phi\rangle \langle \Phi| K_1^{\dagger} = |\Phi'\rangle \langle \Phi'| + |\Phi''\rangle \langle \Phi''|, \qquad (34)$$

where

$$|\Phi'\rangle := a_0 \cos\theta |00\rangle + a_1 \cos\phi |11\rangle, \quad |\Phi''\rangle := a_0 \sin\theta |01\rangle + a_1 \sin\phi |10\rangle.$$
(35)

It is clear that $|\Phi'\rangle \perp |\Phi''\rangle$. Therefore, it results

$$\|\mathrm{id} \otimes \mathcal{N}(\Phi)\|_{1} = \||\Phi'\rangle\langle\Phi'|\|_{1} + \||\Phi''\rangle\langle\Phi''|\|_{1}.$$
(36)

As a consequence we have that Φ allows us to saturate the bound (29), namely

$$\|\Delta_{\mathrm{id}\otimes\mathcal{N}_{\phi,\theta}}(\Phi)\|_{1} = \|\Delta_{\Phi'}\|_{1} + \|\Delta_{\Phi''}\|_{1}.$$
(37)

The same result can be found starting from $|\Phi\rangle = a_0|00\rangle + a_1|11\rangle$.

In the canonical basis the quantities at the l.h.s. of Eq.(37) read:

$$\Delta_{\Phi'} := \begin{pmatrix} |a_0|^2 \alpha & 0 & 0 & a_0 a_1^* \gamma_1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ a_0^* a_1 \gamma_1 & 0 & 0 & |a_1|^2 \beta \end{pmatrix}, \qquad \Delta_{\Phi''} := \begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & -|a_1|^2 \alpha & a_0 a_1^* \gamma_2 & 0 \\ 0 & a_0^* a_1 \gamma_2 & -|a_0|^2 \beta & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}.$$
(38)

Comparing Eq.(38) with Eq.(24) we deduce that

$$\begin{aligned} \|\Delta_{\mathrm{id} \otimes \mathcal{N}_{\phi,\theta}}(\Phi)\|_{1} &= \|\Delta_{\Phi'}\|_{1} + \|\Delta_{\Phi''}\|_{1} \\ &= \|\Delta_{\psi'}\|_{1} + \|\Delta_{\psi''}\|_{1} \ge \|\Delta_{\mathcal{N}_{\phi,\theta}}(\psi)\|_{1}. \end{aligned}$$
(39)

Finally, from Eqs.(37) and (38), we derive

$$\max_{\Psi \in \mathcal{P}(\mathbb{C}^2 \otimes \mathbb{C}^2)} \|\Delta_{\mathrm{id} \otimes \mathcal{N}_{\phi,\theta}}(\Psi)\|_1 = \max_{\Phi \in \mathcal{P}(\mathbb{C}^2)} \|\Delta_{\mathrm{id} \otimes \mathcal{N}_{\phi,\theta}}(\Phi)\|_1$$
$$= \max_{0 \le s \le 1} \frac{1}{2} \sum_{i,j=1}^2 \left| (1-s)\alpha + s\beta + (-1)^i \sqrt{((1-s)\alpha + s\beta)^2 + 4s(1-s)(\gamma_j^2 - \alpha\beta)} \right|.$$
(40)

Observation III.4 First notice that, being $(\gamma_j^2 - \alpha\beta) \ge 0$ for j = 1, 2, Eq.(28) reduces to

$$\max_{\Psi \in \mathcal{P}(\mathbb{C}^2 \otimes \mathbb{C}^2)} \|\Delta_{\mathrm{id} \otimes \mathcal{N}_{\phi,\theta}}(\Psi)\|_1 = \max_{0 \le s \le 1} \left\{ \sqrt{((1-s)\alpha + s\beta)^2 + 4s(1-s)(\gamma_1^2 - \alpha\beta)} + \sqrt{((1-s)\alpha + s\beta)^2 + 4s(1-s)(\gamma_2^2 - \alpha\beta)} \right\}.$$
 (41)

Then, observe that by means of Eqs.(41) and (20), we can write

$$\max_{\Psi \in \mathcal{P}(\mathbb{C}^2 \otimes \mathbb{C}^2)} \|\Delta_{\mathrm{id} \otimes \mathcal{N}_{\phi,\theta}}(\Psi)\|_1 - \max_{\psi \in \mathcal{P}(\mathbb{C}^2)} \|\Delta_{\mathcal{N}_{\phi,\theta}}(\psi)\|_1 = \max_{0 \le s \le 1} f(s) - \max_{0 \le s \le 1} g(s),$$
(42)

where

$$f(s) := \sum_{i=1}^{2} \sqrt{((1-s)\alpha - s\beta)^2 + 4s(1-s)\gamma_i^2},$$
(43)

$$g(s) := 2\sqrt{((1-s)\alpha - s\beta)^2 + 4s(1-s)\left(\frac{|\gamma_1| + |\gamma_2|}{2}\right)^2}.$$
(44)

It follows that:

- When $|\gamma_1| = |\gamma_2|$, it is f(s) = g(s), implying that side entanglement is not useful.
- When $|\gamma_1| \neq |\gamma_2|$, suppose that $\max_{0 \le s \le 1} g(s) = g(s_*)$, then, using relation (39), we will have

$$\max_{0 \le s \le 1} f(s) \ge f(s_*) \ge g(s_*).$$
(45)

In other words

If f (s_{*}) ≠ g(s_{*}) ⇒ max_Ψ ||Δ_{id ⊗N_{φ,θ}}(Ψ)||₁ > max_ψ ||Δ_{N_{φ,θ}}(ψ)||₁ ⇒ side entanglement is useful.
If f (s_{*}) = g(s_{*}) and s_{*} ≠ argmax f (s) ⇒ side entanglement is useful.
If f (s_{*}) = g(s_{*}) and s_{*} = argmax f (s) ⇒ side entanglement is not useful.

With this in mind we can now state the following Theorem.

Theorem III.5 In the discrimination between two qubit channels, $\mathcal{N}_{\phi_1,\theta_1}$ and $\mathcal{N}_{\phi_1,\theta_1}$, that are extreme point maps of \mathfrak{N} , when $|\gamma_1| = |\gamma_2|$ side entanglement is not useful. Instead, when $|\gamma_1| \neq |\gamma_2|$ the following exhausts all possibilities:

- A. Side entanglement is not useful when
 - 1) $2|\gamma_M| \le |\alpha + \beta|;$
 - 2) $|\gamma_1| + |\gamma_2| \le |\alpha + \beta| < 2|\gamma_M|$ and $P(\alpha + \beta) \ge \gamma_1^2 + \gamma_2^2$;
 - 3) $|\alpha + \beta| < |\gamma_1| + |\gamma_2|$ and $\alpha \neq \beta$;
 - 4) $|\alpha + \beta| < |\gamma_1| + |\gamma_2|, \alpha = \beta \text{ and } \alpha^2 \le |\gamma_1| |\gamma_2|.$
- B. Side entanglement is useful when
 - 1) $|\gamma_1| + |\gamma_2| \le |\alpha + \beta| < 2|\gamma_M|$ and $P(\alpha + \beta) < \gamma_1^2 + \gamma_2^2$;
 - 2) $|\alpha + \beta| < |\gamma_1| + |\gamma_2|, \alpha = \beta$, and $\alpha^2 > |\gamma_1| |\gamma_2|$.

Proof When $|\gamma_1| = |\gamma_2|$, the functions we are optimizing become the same, i.e., f(s) = g(s), hence side entanglement cannot be useful as put forward in Observation III.4.

• If $2|\gamma_M| \leq |\alpha + \beta|$, we have from Lemma III.2,

$$Eq.(42) = \max_{0 \le s \le 1} \sum_{i=1}^{2} \sqrt{((1-s)\alpha - s\beta)^{2} + 4s(1-s)\gamma_{i}^{2}} - 2|P|$$

$$\leq \sum_{i=1}^{2} \max_{0 \le s \le 1} \sqrt{((1-s)\alpha - s\beta)^{2} + 4s(1-s)\gamma_{i}^{2}} - 2|P|$$

$$\leq \max\{|\alpha|, |\beta|\} + \max\{|\alpha|, |\beta|\} - 2|P| = 0,$$
(46)

meaning that side entanglement is not useful. This proves A.1).

• When $|\gamma_1| + |\gamma_2| \le |\alpha + \beta| < 2|\gamma_M|$, if there exist a value s, 0 < s < 1, for which $f(s) > g(s_*)$, then side entanglement is useful. Thanks to Lemma III.2 we can rewrite $f(s) > g(s_*)$ as

$$\sum_{i=1}^{2} \sqrt{\left(s\alpha - (1-s)\beta\right)^2 + 4s(1-s)\gamma_i^2} > 2|P|.$$
(47)

Instead, if $\forall s \in (0, 1)$ Eq.(47) does not hold, then side entanglement is not useful.

The two above alternatives are equivalent to the fact that the following function

$$F(s) := P^{2}(P^{2} - \beta^{2}) + 2P^{2} \left(\beta(\alpha + \beta) - (\gamma_{1}^{2} + \gamma_{2}^{2})\right) s + \left[P^{2} \left(2(\gamma_{1}^{2} + \gamma_{2}^{2}) - (\alpha + \beta)^{2}\right) + (\gamma_{1}^{2} - \gamma_{2}^{2})^{2}\right] s^{2} - 2(\gamma_{1}^{2} - \gamma_{2}^{2})^{2} s^{3} + (\gamma_{1}^{2} - \gamma_{2}^{2})^{2} s^{4},$$
(48)

can be negative for some values of *s*, or it is always not negative.

We then distinguish two cases:

i) $P = \alpha$.

- If $\alpha(\alpha + \beta) < \gamma_1^2 + \gamma_2^2$, we can make F(s) < 0 with *s* arbitrarily close to 0. Hence, side entanglement is useful. This partly proves B.1).
- If $\alpha(\alpha + \beta) \ge \gamma_1^2 + \gamma_2^2$, for each 0 < s < 1, we have

$$2\alpha(\alpha+\beta) - (\alpha+\beta)^2(1-s) - 2s(\gamma_1^2+\gamma_2^2) \ge 0,$$
(49)

which implies that $F(s) \ge 0$ and so, side entanglement is not useful. This partly proves A.2).

ii)
$$P = \beta$$
.

- If $\beta(\alpha + \beta) < \gamma_1^2 + \gamma_2^2$, we can make F(s) < 0, with *s* arbitrarily close to 1. Hence side entanglement is useful. This completes the proof of B.1).
- If $\beta(\alpha + \beta) \ge \gamma_1^2 + \gamma_2^2$ for each 0 < s < 1, we have

$$2\beta(\alpha+\beta) - (\alpha+\beta)^2(1-s) - 2s(\gamma_1^2+\gamma_2^2) \ge 0,$$
(50)

which implies that $F(s) \ge 0$ and so, side entanglement is not useful. This completes the proof of A.2).

• When $|\alpha + \beta| < |\gamma_1| + |\gamma_2|$, it is

$$f(s_*) \neq g(s_*) \Leftrightarrow ((1 - s_*)\alpha - s_*\beta) (|\gamma_1| - |\gamma_2|)^2 \neq 0,$$
 (51)

for any $0 < s_* < 1$. Inserting the solution (27) into (51), we can realize that $((1 - s_*)\alpha - s_*\beta)(|\gamma_1| - |\gamma_2|)^2 = 0$ for $\alpha \neq \beta$. Hence, in this case side entanglement is not useful. This proves A.3).

• When $|\alpha + \beta| < |\gamma_1| + |\gamma_2|$ and $\alpha = \beta$, the solution (27) becomes $s_* = \frac{1}{2}$. One can easily check that this is also an extrema of the function f, but it turns out to be the maximum only when $\alpha^2 \le |\gamma_1||\gamma_2|$. It means that if $\alpha^2 > |\gamma_1||\gamma_2|$, side entanglement is useful, proving B.2). Instead if $\alpha^2 \le |\gamma_1||\gamma_2|$ side entanglement is not useful, proving A.4).

Remark The channels of Theorem III.5 are on the faces of the tetrahedron and correspond to ellipsoids that have one point in common with the boundary of the Bloch sphere. Examples of this kind of channels are those having two Kraus operators that are non-unitary, like the amplitude damping channel that can be written as $\mathcal{N}_{\phi,\theta=0}(\rho)$. For it, from Theorem III.5, we can easily recover the results of Ref.[15].

IV. DISCRIMINATING BETWEEN NON-EXTREMAL QUBIT CHANNELS

In this Section we extend the result of Theorem III.5 to the discrimination between two nonextremal qubit channels. These can be written as convex combination of extremal channels as follows:

$$\mathcal{N}_{\{\lambda_1,\phi_1,\theta_1,\phi_1',\theta_1'\}} = \lambda_1 \mathcal{N}_{\phi_1,\theta_1} + (1-\lambda_1) \mathcal{N}_{\phi_1',\theta_1'},\tag{52}$$

$$\mathcal{N}_{\{\lambda_2,\phi_2,\theta_2,\phi'_2,\theta'_2\}} = \lambda_2 \mathcal{N}_{\phi_2,\theta_2} + (1-\lambda_2) \mathcal{N}_{\phi'_2,\theta'_2},\tag{53}$$

where $0 < \lambda_1, \lambda_2 < 1$.

1

Then, we redefine the parameters of (14), (15), (16) and (17) in a more general fashion as

$$\alpha := (\lambda_1 \cos^2 \theta_1 + (1 - \lambda_1) \cos^2 \theta_1') - (\lambda_2 \cos^2 \theta_2 + (1 - \lambda_2) \cos^2 \theta_2'), \tag{54}$$

$$\beta := (\lambda_1 \cos^2 \phi_1 + (1 - \lambda_1) \cos^2 \phi_1') - (\lambda_2 \cos^2 \phi_2 + (1 - \lambda_2) \cos^2 \phi_2'), \tag{55}$$

$$\gamma_1 := \lambda_1 \cos \phi_1 \cos \theta_1 + (1 - \lambda_1) \cos \phi_1' \cos \theta_1' - \lambda_2 \cos \phi_2 \cos \theta_2 - (1 - \lambda_2) \cos \phi_2' \cos \theta_2', \tag{56}$$

$$\gamma_2 \coloneqq \lambda_1 \sin \phi_1 \sin \theta_1 + (1 - \lambda_1) \sin \phi_1' \sin \theta_1' - \lambda_2 \sin \phi_2 \sin \theta_2 - (1 - \lambda_2) \sin \phi_2' \sin \theta_2'. \tag{57}$$

Remark Note that all the machinery of Lemmas III.2, III.3 and Observation III.4 applies to parameters (54)-(57) as well. At the core of this is the fact that, proceeding similarly to Eq.(34), we can now find

$$\mathrm{id} \otimes \mathcal{N}_{\{\lambda,\phi,\theta,\phi',\theta'\}}(\Phi) = \lambda_1 \left(|\Phi'\rangle \langle \Phi'| + |\Phi''\rangle \langle \Phi''| \right) + (1 - \lambda_1) \left(|\Gamma'\rangle \langle \Gamma'| + |\Gamma''\rangle \langle \Gamma''| \right), \tag{58}$$

where $|\Phi'\rangle$, $|\Phi''\rangle$ are as in Eq.(35), while $|\Gamma'\rangle := a_0 \cos \theta'_1 |00\rangle + a_1 \cos \phi'_1 |11\rangle$, $|\Gamma''\rangle := a_0 \sin \theta'_1 |01\rangle + a_1 \sin \phi'_1 |10\rangle$. Then, since $\{|\Phi'\rangle, |\Gamma'\rangle\} \perp \{|\Phi''\rangle, |\Gamma''\rangle\}$, we can draw a conclusion analogous to Eq.(29), namely

$$\|\Delta_{\mathrm{id}\otimes\mathcal{N}}(\Phi)\|_{1} = \|\lambda_{1}\Delta_{\Phi'} + (1-\lambda_{1})\Delta_{\Gamma'}\|_{1} + \|\lambda_{1}\Delta_{\Phi''} + (1-\lambda_{1})\Delta_{\Gamma''}\|_{1}.$$
(59)

However, parameters (54)-(57) no longer guarantee the satisfaction of conditions $(\gamma_j^2 - \alpha\beta) \ge 0$, for j = 1, 2, as in Sec.III. Whenever $\gamma_m^2 \ge \alpha\beta$, this holds true and Theorem III.5 is directly applicable. For the other cases we have the following results.

Theorem IV.1 In the discrimination between two qubit channels, $\mathcal{N}_{\{\lambda_1,\phi_1,\theta_1,\phi_1',\theta_1'\}}$ and $\mathcal{N}_{\{\lambda_2,\phi_2,\theta_2,\phi_2',\theta_2'\}}$, that are not extreme point maps of \mathfrak{N} , when $\gamma_M^2 \leq \alpha\beta$ side entanglement is not useful. Instead, when $\gamma_m^2 < \alpha\beta < \gamma_M^2$ the following exhausts all possibilities:

- A. Side entanglement is not useful when
 - 1) $|\gamma_1| + |\gamma_2| > |\alpha + \beta|$, while $\alpha = \beta$ and $|\gamma_m| = |\alpha|$;
 - 2) $|\gamma_1|+|\gamma_2| > |\alpha+\beta|$ while $\alpha \neq \beta$, $2\gamma_M(\alpha+\beta) = (|\gamma_1|+|\gamma_2|)^2$, and $\gamma_M = \frac{1+\alpha\beta\pm\sqrt{(\alpha^2-1)(\beta^2-1)}}{\alpha+\beta}$;
 - 3) $|\gamma_1| + |\gamma_2| \le |\alpha + \beta|$ and $|\gamma_M| \le |P|$.

B. Side entanglement is useful when

- 1) $|\gamma_1| + |\gamma_2| > |\alpha + \beta|$, while $\alpha = \beta$ and $|\gamma_m| \neq |\alpha|$;
- 2) $|\gamma_1| + |\gamma_2| > |\alpha + \beta|$, while $\alpha \neq \beta$ and $2\gamma_M(\alpha + \beta) \neq (|\gamma_1| + |\gamma_2|)^2$;
- 3) $|\gamma_1| + |\gamma_2| > |\alpha + \beta|$, while $\alpha \neq \beta$, $2\gamma_M(\alpha + \beta) = (|\gamma_1| + |\gamma_2|)^2$, and $\gamma_M \neq \frac{1 + \alpha\beta \pm \sqrt{(\alpha^2 1)(\beta^2 1)}}{\alpha + \beta}$; 4) $|\gamma_1| + |\gamma_2| \le |\alpha + \beta|$ and $|\gamma_M| > |P|$.

Proof If $\gamma_m^2 \leq \alpha\beta$, on the one hand we have from Eq.(28) $\max_{\Psi} \|\Delta_{\operatorname{id} \otimes \mathcal{N}}(\Psi)\|_1 = \max_{0 \leq s \leq 1} |(1 - s)\alpha + s\beta|$, on the other hand from Lemma III.2 we have $\max_{\psi} \|\Delta_{\mathcal{N}}(\psi)\|_1 \geq \max_{0 \leq s \leq 1} |(1 - s)\alpha + s\beta|$. Thus

$$\max_{\Psi} \|\Delta_{\mathrm{id}\,\otimes\mathcal{N}}(\Psi)\|_{1} - \max_{\psi} \|\Delta_{\mathcal{N}}(\psi)\|_{1} \le \max_{0\le s\le 1} |(1-s)\alpha + s\beta| - \max_{0\le s\le 1} |(1-s)\alpha + s\beta| = 0, \quad (60)$$

meaning that side entanglement is not useful.

The rest of the proof is devoted to the case $\gamma_m^2 < \alpha\beta < \gamma_M^2$. In such a case Eq.(28) reduces to

$$\max_{\Psi} \|\Delta_{\mathrm{id}\otimes\mathcal{N}}(\Psi)\|_{1} = \max_{0\leq s\leq 1} \left(|(1-s)\alpha + s\beta| + \sqrt{((1-s)\alpha + s\beta)^{2} + 4s(1-s)(\gamma_{M}^{2} - \alpha\beta)} \right).$$
(61)

Let us define

$$G(s) := |(1-s)\alpha + s\beta| + \sqrt{((1-s)\alpha + s\beta)^2 + 4s(1-s)(\gamma_M^2 - \alpha\beta)}.$$
(62)

This function has maxima for the argument $s_* = \min\{\max\{0, \tilde{s}\}, 1\}$, where

$$\tilde{s} = \begin{cases} \frac{\gamma_M - \alpha}{2\gamma_M - (\alpha + \beta)}, & \text{when } \gamma_M(\alpha + \beta) > 0\\ \\ \frac{\gamma_M + \alpha}{2\gamma_M + (\alpha + \beta)}, & \text{when } \gamma_M(\alpha + \beta) < 0 \end{cases},$$
(63)

leading to

$$\max_{\Psi} \|\Delta_{\mathrm{id} \otimes \mathcal{N}}(\Psi)\|_{1} = \begin{cases} |\gamma_{M}| + \left|\frac{2\alpha\beta - \gamma_{M}(\alpha+\beta)}{\alpha+\beta-2\gamma_{M}}\right| & \text{when} \quad \gamma_{M}(\alpha+\beta) > 0\\ & & \\ |\gamma_{M}| + \left|\frac{2\alpha\beta + \gamma_{M}(\alpha+\beta)}{\alpha+\beta+2\gamma_{M}}\right| & \text{when} \quad \gamma_{M}(\alpha+\beta) < 0 \end{cases}$$
(64)

• Consider now $|\gamma_1| + |\gamma_2| > |\alpha + \beta|$. We distinguish the following cases:

- 1. When $\alpha \neq \beta$ and $2\gamma_M(\alpha + \beta) \neq (|\gamma_1| + |\gamma_2|)^2$, using (27) and (63), we have $s_* \neq |a_1|^2$, meaning that side entanglement is useful. This proves B.2).
- 2. When $\alpha \neq \beta$ and $2\gamma_M(\alpha + \beta) = (|\gamma_1| + |\gamma_2|)^2$ we have $\gamma_M(\alpha + \beta) > 0$, and $s^* = |a_1|^2 = \frac{\gamma_M + \alpha}{2\gamma_M + (\alpha + \beta)}$ which yields

$$\frac{\max_{\Psi} \|\Delta_{\mathrm{id}\otimes\mathcal{N}}(\Psi)\|_{1}}{\max_{\psi} \|\Delta_{\mathcal{N}}(\psi)\|_{1}} = \frac{1}{2} \left(\sqrt{\frac{\alpha + \beta - 2\gamma_{M}}{\gamma_{M} \left(2\alpha\beta - \gamma_{M}(\alpha + \beta)\right)}} + \sqrt{\frac{\gamma_{M} \left(2\alpha\beta - \gamma_{M}(\alpha + \beta)\right)}{\alpha + \beta - 2\gamma_{M}}} \right).$$
(65)

Side entanglement is useful when (65)> 1. This happens for $\gamma_M \neq \frac{1+\alpha\beta\pm\sqrt{(\alpha^2-1)(\beta^2-1)}}{\alpha+\beta}$, proving B.4). In contrast, for $\gamma_M = \frac{1+\alpha\beta\pm\sqrt{(\alpha^2-1)(\beta^2-1)}}{\alpha+\beta}$ it results (65)=1 and side entanglement is useless. This proves A.2).

3. When $\alpha = \beta$, again using (27) and (63), we have $s_* = |a_1|^2 = \frac{1}{2}$ leading to

$$G\left(\frac{1}{2}\right) = |\gamma_M| + |\alpha| \quad \text{and} \quad g\left(\frac{1}{2}\right) = |\gamma_M| + |\gamma_m|,$$
(66)

with G, g defined respectively in (62), (44). Thus, we have that

- for $|\gamma_m| \neq |\alpha|$ it is $G(1/2) \neq g(1/2)$ and side entanglement is useful. This proves B.1);
- for $|\gamma_m| = |\alpha|$ it is G(1/2) = g(1/2) and side entanglement is not useful. This proves A.1).

• Finally, when $|\gamma_1| + |\gamma_2| \le |\alpha + \beta|$, if there exist a value s, 0 < s < 1 for which $G(s) > g(s_*)$, then side entanglement is useful. Actually we can rewrite $G(s) > g(s_*)$ as

$$|(1-s)\alpha + s\beta| + \sqrt{((1-s)\alpha - s\beta)^2 + 4(s-s^2)\gamma_M^2} > 2|P|.$$
(67)

Instead, if $\forall s \in (0, 1)$ Eq.(67) does not hold, then side entanglement is not useful.

The two above alternatives are equivalent to the fact that the following function

$$R(s) := P^{2}(P^{2} - \alpha^{2}) + 2P^{2} \left(\alpha^{2} - \gamma_{M}^{2}\right) s + \left[(\gamma_{M}^{2} - \alpha\beta)^{2} - P^{2} \left(\alpha^{2} + \beta^{2} - 2\gamma_{M}^{2}\right)\right] s^{2} - 2s^{3} (\gamma_{M}^{2} - \alpha\beta)^{2} + (\gamma_{M}^{2} - \alpha\beta)^{2} s^{4},$$
(68)

can be negative for some values of *s*, or it is always not negative.

- For $|\gamma_M| > |P|$, we can choose *s* arbitrarily close to zero (or to 1) to make R(s) < 0. This shows that side entanglement is useful and proves B.5).
- For $|\gamma_M| \leq |P|$, for any $0 \leq s \leq 1$, we have R(s) > 0. It means that side entanglement is not useful, proving A.3).

Remark The channels of Theorem IV.1 are inside the tetrahedron and correspond to ellipsoids that have no point in common with the boundary of the Bloch sphere. Examples of this kind of channels are the quantum Pauli channels, that can be written as $\lambda N_{\theta,\theta}(\rho) + (1-\lambda)N_{\theta',\pi+\theta'}(\rho)$. For them, from Theorem IV.1, we can easily recover the results of Ref.[3].

V. CONCLUSION

In summary, we investigated the conditions under which side entanglement can effectively be used for channel discrimination and identified the specific circumstances in which it enhances, as well as those in which it does not enhance, the performance of discerning between two generic qubit channels, up to unitary pre- and post-processing. We present our findings in a decision tree, which is included in Appendix A.

Our results were obtained through a constructive approach, whereby we first established the uselessness of side entanglement in discriminating quasi-extremal channels within the set \mathfrak{N} . We then analyzed the problem for extremal maps and finally extended our findings to generic channels that may be located within such a set.

Our results can be straightforwardly extended to channels that include pre and post-processing unitaries, provided that these are diagonal or anti-diagonal in the canonical basis. This is because, such unitaries will preserve the two subspaces where the optimal input lives ($span\{|00\rangle, |11\rangle\}$, or $span\{|01\rangle, |10\rangle\}$) or map one into another. Furthermore, relations (24-25) and (38-39) will still hold true for newly defined parameters α, β, γ_1 and γ_2 .

Several partial results previously obtained for specific kind of qubit channels can be recovered as particular cases of our general results. Examples are the discrimination of dephasing channel, amplitude damping channel and depolarizing channel.

From Theorems III.5 and IV.1 appears that, quite generally, side entanglement is useful only in a limited portion of the parameters' space. Additionally, the profit is often brought by nonmaximally entangled states.

Our findings are based on the characterization of qubit channels in terms of affine maps on \mathbb{R}^3 . Geometric characterization of the action of quantum channels in higher dimensional systems would be useful for extending our results beyond qubit channels.

Moving on from the framework of point-to-point quantum channel [17], one could shift to more complex structures to discriminate, namely quantum networks [18]. In contrast to channels, networks allow for intermediate access points where information can be received, processed and reintroduced into the network. A first attempt in this direction has been done in Ref. [19] using the idea of quantum super-channels.

Appendix A



FIG. 1. Decision tree for the utility of side entanglement in qubit channel discrimination. Rectangles stand for tests and ellipses for decisions. Symbols refer to Eqs.(54)-(57) with γ_m , γ_M and P given in (18), (19).

Acknowledgments

S.M. and M.R. acknowledge funding from the European Union's Horizon 2020 research and innovation programme, under grant agreement QUARTET No 862644.

^[1] A. Acin, *Statistical distinguishability between unitary operations*, Physical Review Letters, **87**, 177901 (2001).

^[2] A. Gilchrist, N. K. Langford, and M. A. Nielsen, *Distance measures to compare real and ideal quantum processes*, Physical Review A **71**, 062310 (2005).

^[3] M. Sacchi, Optimal discrimination of quantum operations, Physical Review A 71, 062340 (2005).

- [4] G. Wang, and M. Ying, Unambiguous discrimination among quantum operations, Physical Review A 73, 042301 (2006).
- [5] S. Pirandola, C. Lupo, V. Giovannetti, S. Mancini, and S. L. Braunstein, *Quantum reading capacity*, New Journal of Physics 13, 113012 (2011).
- [6] A. M. Childs, J. Preskill, and J. Renes, *Quantum information and precision measurement*, Journal of Modern Optics, 47, 155-176 (2000).
- [7] B. Rosgen, and J. Watrous, On the hardness of distinguishing mixed-state quantum computations, Proceedings of the 20th IEEE Conference on Computational Complexity, pp. 344-354 (2005).
- [8] C.W. Helstrom, Quantum Detection and Estimation Theory, Academic Press, New York (1976).
- [9] C. Datta, T.Biswas, D. Saha, and R. Augusiak, *Perfect discrimination of quantum measurements using entangled systems*, New Journal of Physics 23, 043021(2021).
- [10] M. Sacchi, Entanglement can enhance the distinguishability of entanglement-breaking channels, Physical Review A 72, 014305 (2005).
- [11] D. Puzzuoli, and J. Watrous, Ancilla Dimension in Quantum Channel Discrimination, Annales Henri Poincaré 18, 1153-1184 (2017).
- [12] M. B. Ruskai, S. Szarek, and E. Werner, An analysis of completely positive trace-preserving maps on M₂, Linear Algebra and its Applications 347, 159-187 (2002).
- [13] F. Caruso, V. Giovannetti, C. Lupo and S. Mancini, *Quantum channels and memory effects*, Reviews of Modern Physics 86, 1203 (2014).
- [14] M. Rexiti, L. Memarzadeh, and S. Mancini, *Discrimination of dephasing channels*, Journal of Physics A: Mathematical and Theoretical 55, 245301 (2022).
- [15] M. Rexiti, and S. Mancini, *Discriminating qubit amplitude damping channels*, Journal of Physics A: Mathematical and Theoretical 54, 165303 (2021).
- [16] M. F. Sacchi, *Minimum error discrimination of Pauli channels*, Journal of Optics B: Quantum and Semiclassical Optics 7, S333 (2005).
- [17] L. Gyongyosi, S. Imre, and H.V. Nguyen, A Survey on Quantum Channel Capacities, IEEE Communications Surveys and Tutorials 20, 1149-1205 (2018).
- [18] L. Gyongyosi, and S. Imre, *Advances in the Quantum Internet*, Communications of the ACM **65**, 52-63 (2022).
- [19] C. Hirche, Quantum Network Discrimination, arXiv:2103.02404.
- [20] Analogously entanglement can help in distinguishing between quantum measurements [9].
- [21] A single point has to be interpreted as a pair of degenerate points in some special cases.
- [22] Actually ω_1 and ω_2 can also be two parameters' vectors.